

SCALA-GOPHER

CSP WITH IDIOMATIC SCALA

<https://github.com/rssh/scala-gopher>

goo.gl/dbT3P7

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KharkivPY Group: Scala & Fun, 2016

2. Outline

- ❖ Theory & History
 - ❖ CSP = Communication Sequence Processes
 - ❖ π -calculus (CCS = Calculus of Communicating System)
 - ❖ Languages: (Occam,Limbo, Go[Clojure, Scala, ...])
- ❖ Main constructions / idioms, how they looks
 - ❖ Channels, Selectors, Transputers
 - ❖ Implementation Techniques

3. Theory & History

- ❖ If you familiar with basic concepts, skip to slide 8
- ❖ Just show me scala API: skip to slide 14

4. Theory & History

- ❖ CPS = Communication Sequence Processes.
- ❖ CSS = Calculus of Communicating System.
 - ❖ 1978 First CSP Paper by Tony Hoar.
 - ❖ 1980. CSS by Robert Milner. (Algebraic Processes)
 - ❖ 1985 CSP formalism (influenced by CSS) in CSP Book
 - ❖ <http://www.usingcsp.com/>
 - ❖ 1992 π -calculus [CSS + Channels] (Robert Milner, Joachim Parrow, David Walker)
 - ❖ (large family of Process Algebras, including Actor Model, ACP,)

5. CSP Notation(basic)

- ❖ (A, B, C ...) — processes,
- ❖ (x, y, z ...) — events
 - ❖ atomic: x , STOP, BEEP, or c.v (channel/value)
 - ❖ $c!v$ - send v to channel c (after this c.v happens)
 - ❖ $c?v$ - receive v from channel c (wait until c.v)
- ❖ trace(...) — sequence of events.

6. CSP Notation(basic)

Operations on processes:

$a \rightarrow P$ // after event

$P ||| Q$ // interleave concurrently $P; Q$ // seq

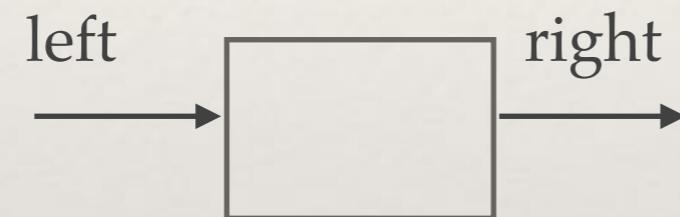
$(a \rightarrow P) \square (b \rightarrow Q)$ // choice, if a then P if b then Q

$\mu X. F(X) \equiv F(\mu X. F(X))$ // fix point (loops, recursion)

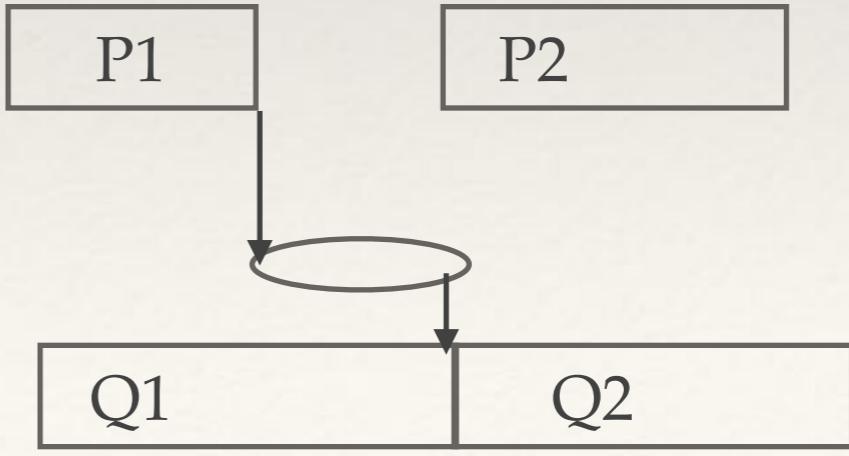
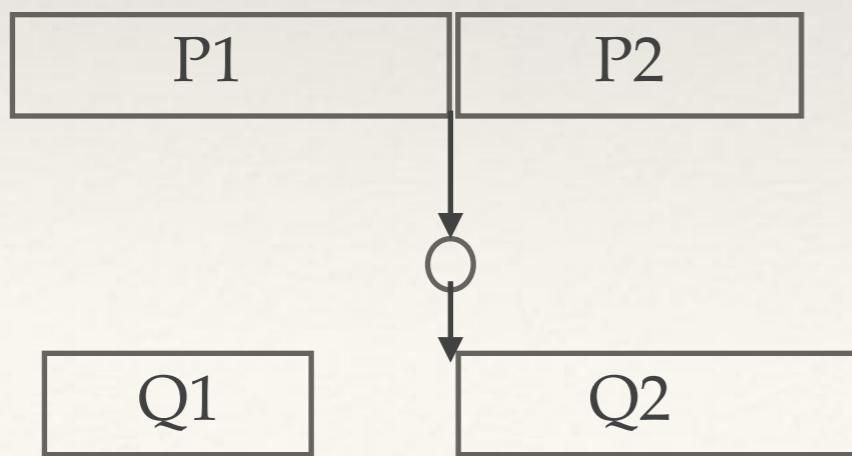
7. CSP Notation(examples)

COPY(left,right)=

$$\mu F. \text{left?}x \rightarrow \text{right!}x \rightarrow F$$



$$(P_1 \rightarrow c!x \rightarrow P_2) ||| (Q_1 \rightarrow c?x \rightarrow Q_2(x))$$



8. Occam language



William Occam, 1287-1347
'Occam razor' principle

Occam language. (1983)

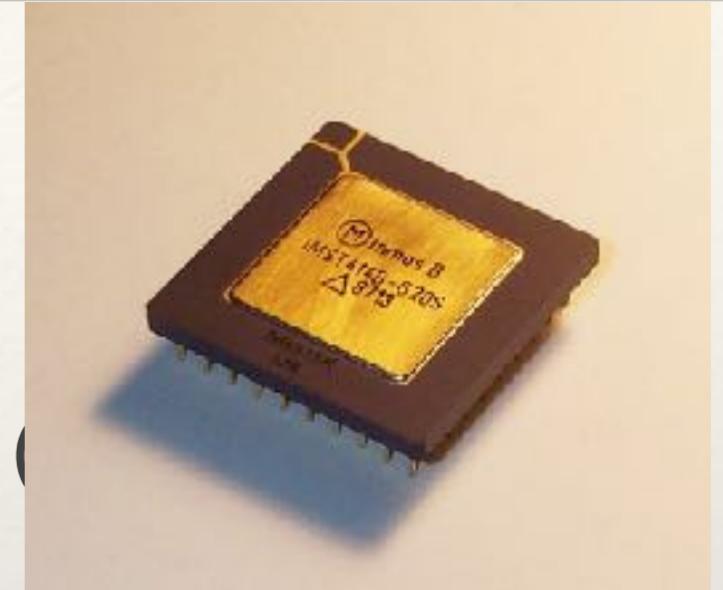
```
INT x, y:  
SEQ  
  x := 4  
  y := (x + 1)  
CHAN INT c:  
PAR  
  some.procedure (x, y, c!)  
  another.procedure (c?)  
  y := 5
```

minimal language. (processor constructors)

braces via indentation. (before python)

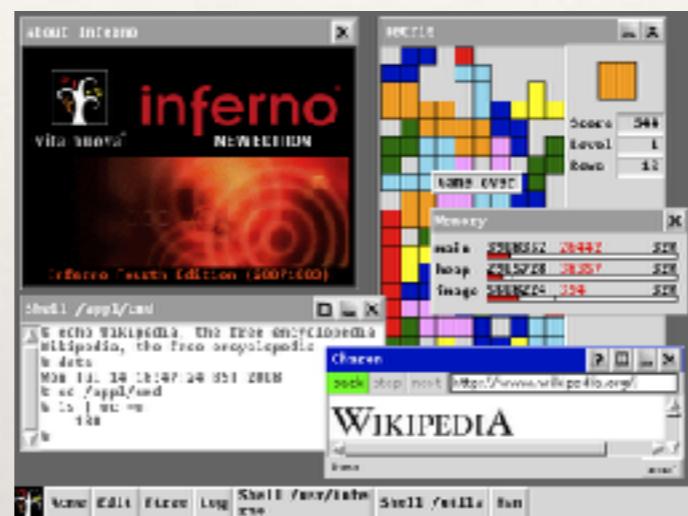
9. Occam language

- ❖ created by INMOS
- ❖ targeted Transputers CHIP architecture
- ❖ latest dialect: occam- π [1996]
 - ❖ extensions from π -calculus
 - ❖ (more dynamic, channels can be send via channels)
- ❖ <http://concurrency.cc> — occam for Aurdino.



10. Inferno, Limbo, Go

- ❖ Unix, Plan9, Inferno: [At&t; Bell labs; vita nuova]



Sean Forward
David Leo Presotto
Rob Pike
Dennis M. Ritchie
Ken Thompson

- ❖ <http://www.vitanuova.com/inferno/>
- ❖ Limbo ... C-like language + channels + buffered channels.
- ❖ Go ... channels as in Limbo (Go roots is from hell is not a metaphor)

11. CPS in Limbo/Go/Scala: main constructions

- ❖ processes: operator for spawning new lightweight process.
- ❖ channels: can be unbuffered (synchronised) or buffered
 - ❖ unbuffered — writer wait until reader start to work
 - ❖ buffered — if channel buffer is not full, writer not blocked
- ❖ selector: wait for few events (channel), eval first.

12. Go: simple code example

Go

```
func fibonacci(c chan int, quit chan bool) {  
    go {  
        x, y := 0, 1  
        for (){  
            select {  
                case c <- x :  
                    x, y = y, x+y  
                case q<-quit:  
                    break;  
            }  
        }  
        close(c)  
    }  
}
```

```
c = make(chan int);  
quit= make(chan bool);  
fibonacci(c,quit)  
for i := range c {  
    fmt.Println(i)  
    if (i > 2000) {  
        quit <- true  
    }  
}
```

13. Go: simple code example

Go

```
func fibonacci(c chan int, quit chan bool) {  
    go {  
        x, y := 0, 1  
        for (){  
            select {  
                case c <- x :  
                    x, y = y, x+y  
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                    break;  
            }  
        }  
        close(c)  
    }  
}
```

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c = make(chan int);  
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fibonacci(c,quit)  
for i := range c {  
    fmt.Println(i)  
    if (i > 2000) {  
        quit <- true  
    }  
}
```

scala-gopher

- ❖ Akka extension + Macros on top of SIP22-async
- ❖ Integrate CSP Algebra and scala concurrency primitives
- ❖ Provides:
 - ❖ asynchronous API inside general control-flow
 - ❖ pseudo-synchronous API inside go{ .. } or async{ .. } blocks
- ❖ Techreport: goo.gl/dbT3P7

```
def nPrimes(n:Int):Future[List[Int]]= {
    val in = makeChannel[Int]()
    val out = makeChannel[Int]()
    go {
        for(i <- 1 to Int.MaxValue) in.write(i)
    }
    go {
        select.fold(in){ (ch,s) =>
            s match {
                case p:ch.read => out.write(p)
                    ch.filter(_ % p != 0)
            }
        }
    }
    go {
        for(i <- 1 to n) yield out.read
    }
}
```

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def nPrimes(n:Int):Future[List[Int]]= {
    val in = makeChannel[Int]()
    val out = makeChannel[Int]()
    go {
        for(i <- 1 to Int.MaxValue)
            in.write(i)
    }
    s match {
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                        ch.filter(_ % p != 0)
                }
            }
        }
        for(i <- 1 to n) yield out.read
    }
}
```

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        select.fold(in){ (ch,s) =>
            s match {
                case p:ch.read => out.write(p)
                    ch.filter(_ % p != 0)
            }
        }
    }
    go {
        for(i <- 1 to n) yield out.read
    }
}
```

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    }
    go {
        select.fold(in){ (ch,s) =>
            s match {
                case p:ch.read => out.write(p)
                    ch.filter(_ % p != 0)
            }
        }
    }
    go {
        for(i <- 1 to n) yield out.read
    }
}
```

Goroutines

- ❖ `go[X](body: X):Future[X]`
 - ❖ Wrapper around `async` +
 - ❖ translation of high-order functions into `async` form
 - ❖ handling of `defer` statement

Goroutines

- ❖ translation of hight-order functions into async form
 - ❖ $f(g)$: $f: (A \Rightarrow B) \Rightarrow C$ in $g: A \Rightarrow B$,
 - ❖ g is invocation-only in f iff
 - ❖ g called in f or in some h inside $f : g$ invocation-only in h
 - ❖ g is
 - ❖ not stored in memory behind f
 - ❖ not returned from f as return value
 - ❖ Collection API high-order methods are invocation-only

Translation of invocation-only functions

- ❖ $f: ((A \Rightarrow B) \Rightarrow C)$, $g: (A \Rightarrow B)$, g invocation-only in f
- ❖ $f': ((A \Rightarrow \text{Future}[B]) \Rightarrow \text{Future}[C])$ $g': (A \Rightarrow \text{Future}[B])$
 - ❖ $\text{await}(g') == g \Rightarrow \text{await}(f') == f$
 - ❖ $f' \Rightarrow \text{await}[\text{translate}(f)]$
 - ❖ $g(x) \Rightarrow \text{await}(g'(x))$
 - ❖ $h(g) \Rightarrow \text{await}(h'(g'))$ iff g is invocation-only in h
 - ❖ That's all
 - ❖ (implemented for specific shapes and parts of scala collection API)

```
def nPrimes(n:Int):Future[List[Int]]= {
    val in = makeChannel[Int]()
    val out = makeChannel[Int]()
    go {
        for(i <- 1 to n*n) in.write(i)
    }
    go {
        select.fold(in){ (ch,s) =>
            s match {
                case p:ch.read => out.write(p)
                    ch.filter(_ % p != 0)
            }
        }
    }
    go {
        for(i <- 1 to n) yield out.read
    }
}
```

```
go {
    (1 to n).map(i => out.read)
}
```

```
async{
    await(t[(1 to n).map(i => out.read)])
}
```

```
async{
    await((1 to n).mapAsync(t[i => async(out.read)]))
}
```

```
async{
    await((1 to n).mapAsync(i => async(await(out.aread))))
}
```

```
mapAsync(i => out.aread)
```

Channels

- ❖ Channel[A] <: Input[A] + Output[A]
 - ❖ Unbuffered
 - ❖ Buffered
 - ❖ Dynamically growing buffers [a-la actor mailbox]
 - ❖ One-time channels [Underlaying promise/Future]
 - ❖ Custom
- 
- CSP

Input[A] - internal API

```
trait Input[A]
{
    type read = A
    // ContRead[A,B].F
    def cbread(f: ContRead[A,B]=>Option[
        ContRead.In[A] => Future[Continuated[B]])
```

```
case class ContRead[A,B](
    function: F,
    channel: Channel[A],
    flowTermination: FlowTermination[B]
)
```

```
// in ContRead companion object
sealed trait In[+A]
case class Value(a:A) extends In[A]
case class Failure(ex: Throwable) extends In[Nothing]
case object Skip extends In[Nothing]
case object ChannelClosed extends In[Nothing]
```

- Continuated[B]
- ContRead
- ContWrite
- Skip
- Done
- Never

Input[A] - external API

```
trait Input[A]  
{
```

```
.....
```

```
def aread: Future[A] = <implementation...>
```

```
def read: A = macro <implementation ... >
```

```
await(aread)
```

```
....
```

```
def map[B](f: A=>B): Input[B] = ....
```

```
// or, zip, filter, ... etc
```

+ usual operations on streams in functional language

Output[A] - API

```
trait Output[A]  
{
```

```
    type write = A
```

ContWrite[A,B].F



```
    def cbwrite(f: ContWrite[A,B]=>Option[  
        (A, Future[Continuated[B]])),  
        ft: FlowTermination[B])  
.....
```

```
    def awrite(a:A): Future[A] = ....
```

```
    def write(a:A): A = .... << await(awrite)
```

```
....
```

```
case class ContWrite[A,B](  
    function: F,  
    channel: Channel[A],  
    flowTermination: FlowTermination[B]  
)
```

Selector

$$(a \rightarrow P) \square (b \rightarrow Q)$$

Go language:

```
go {
  for{
    select{
      case c1 -> x: ... // P
      case c2 <- y: ... // Q
    }
  }
}
```

$*[(c_1?x \rightarrow P) \square (c_2!y \rightarrow Q)]$

Scala:

```
go {
  select.forever {
    case x : c1.read => ...
    case y : c2.write => ...
  }
}
```

```
select.aforever {
  case x : c1.read => ...
  case y : c2.write => ...
}
```

Provide set of flow combinators:
forever, once, fold

select: fold API

```
def fibonacci(c: Output[Long], quit: Input[Boolean]): Future[(Long,Long)] =  
  select.afold((0L,1L)) { case ((x,y),s) =>  
    s match {  
      case x: c.write => (y, x+y)  
      case q: quit.read =>  
        select.exit((x,y))  
    }  
  }
```

fold/afold:

- special syntax for tuple support
- 's': selector pseudoobject
- s match must be the first statement
- select.exit(..) to return value from flow

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            }
        }
    }
    go {
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await(aread)
```

```
....
```

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// or, zip, filter, ... etc
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+ usual operations on streams in functional language

Output[A] - API

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trait Output[A]  
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```
    type write = A
```

ContWrite[A,B].F



```
    def cbwrite(f: ContWrite[A,B]=>Option[  
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```

```
    def awrite(a:A): Future[A] = ....
```

```
    def write(a:A): A = .... << await(awrite)
```

```
....
```

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Selector

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Go language:

```
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  for{
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$*[(c_1?x \rightarrow P) \square (c_2!y \rightarrow Q)]$

Scala:

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select.aforever {
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  case y : c2.write => ...
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Provide set of flow combinators:
forever, once, fold

select: fold API

```
def fibonacci(c: Output[Long], quit: Input[Boolean]): Future[(Long,Long)] =  
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fold/afold:

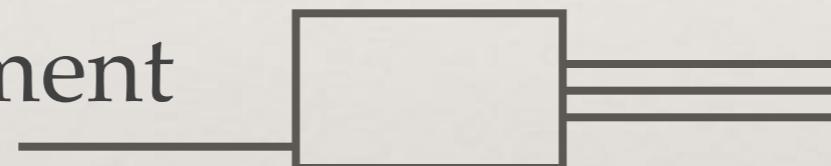
- special syntax for tuple support
- 's': selector pseudoobject
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Transputer



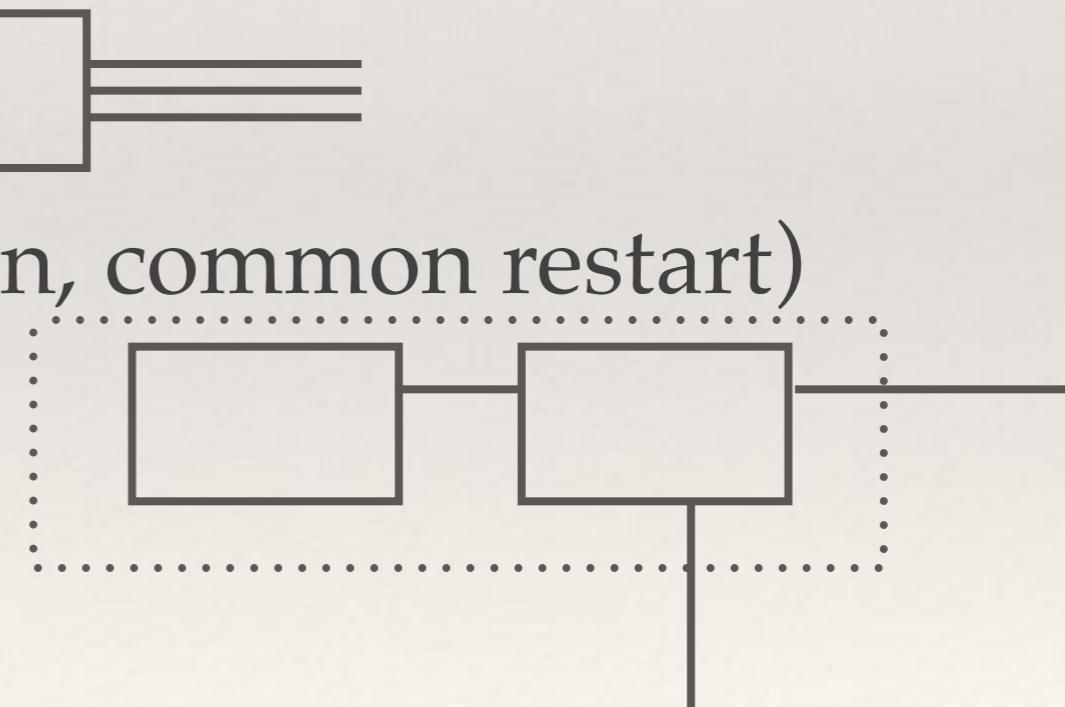
- ❖ Actor-like object with set of input / output ports, which can be connected by channels
- ❖ Participate in actor systems supervisors hierarchy

- ❖ SelectStatement



- ❖ A+B (parallel execution, common restart)

- ❖ replicate



Channel-based generic API

```
val listeners: Channel[Channel[T]]  
val messages: Channel[T] = makeChannel[]
```

```
// private part  
case class Message(next:Channel[Message],value:T)
```

```
select.afold(makeChannel[Message]) { (bus, s) =>  
    s match {  
        case v: message.read => val newBus = makeChannel[Message]  
            current.write(Message(newBus,v))  
            newBus  
        case ch: listener.read => select.afold(bus) { (current,s) =>  
            s match {  
                case msg:current.read => ch.awrite(msg.value)  
                    current.write(msg)  
                    msg.next  
            }  
        }  
    }  
}
```

- state - channel [bus], for which all listeners are subscribed
 - on new message - send one to bus with pointer to the next bus state
 - listener on new message in bus - handle, change current and send again
 - on new listener - propagate

Channel-based generic API

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val listener: Channel[Channel[T]]  
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            s match {  
                case msg:current.read => ch.awrite(msg.value)  
                    current.write(msg)  
                    msg.next  
            }  
            current  
        }  
    }  
}
```

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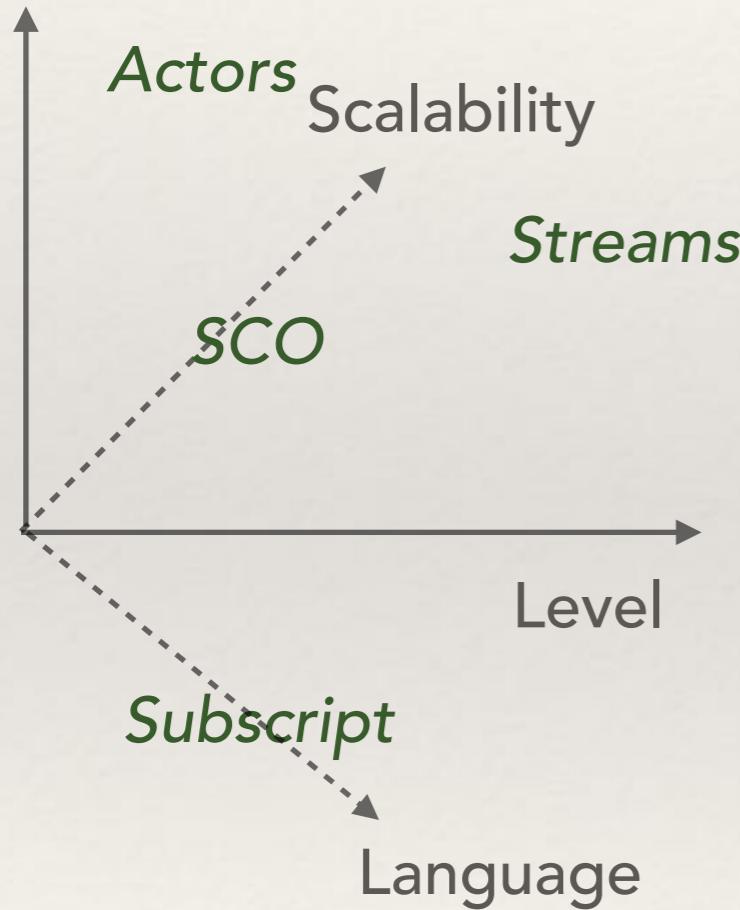
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            s match {  
                case msg:current.read => ch.awrite(msg.value)  
                                            current.write(msg)  
                                            msg.next  
            }  
        }  
        current
```

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Scala concurrency libraries

Flexibility

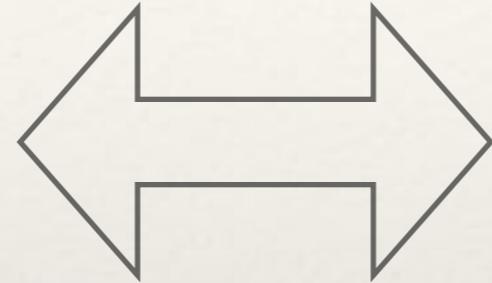


- **Actors**
 - *low level,*
 - *great flexibility and scalability*
- **Akka-Streams**
 - *low flexibility*
 - *high-level, scalable*
- **SCO**
 - *low scalability*
 - *high-level, flexible*
- **Reactive Isolated**
 - *high-level, scalable,*
 - *allows delegation*
- **Gopher**
 - *can emulate each style*

Gopher vs Reactive Isolates

Gopher

- Transputer/fold
- Input
- Output



Isolates

- Isolate
- Events
- Channel

Many writers

One writer

Channel must have owner

CSP + growing buffer

Loosely coupled (growing buffer)

Local

Distributed

Scala-gopher: early experience reports

- ❖ Not 1.0 yet
- ❖ Helper functionality in industrial software projects.
(utilities, small team)
- ❖ Generally: positive
 - ❖ transformation of invocation-only hight-order methods into async form
 - ❖ recursive dynamic data flows
- ❖ Error handling needs some boilerplate

Error handling: language level issue

```
val future = go {  
    .....  
    throw some exception  
}
```

```
go {  
    .....  
    throw some exception  
}
```

```
Go {  
    .....  
    throw some exception  
}
```

Core scala library:
Future.apply
(same issue)

Error is ignored

Developers miss-up Go/go

Errors in ignored value: possible language changes.

- ❖ Possible solutions:
 - ❖ Optional implicit conversion for ignored value
 - ❖ Special optional method name for calling with ignored value
 - ❖ Special return type

```
trait Ignored[F]
```

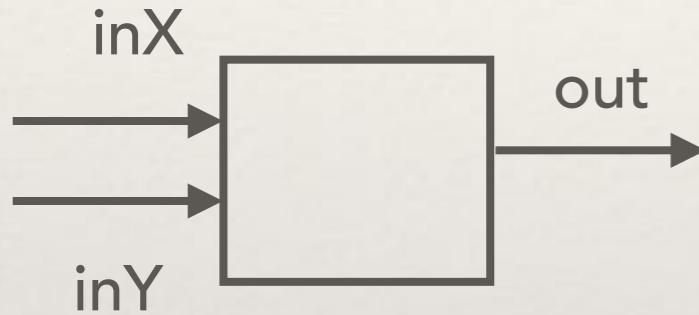
```
object Future
{
    implicit def toIgnored(f:Future):Ignored[Future] =
    ....
```

```
def go[X](f: X): Future[X]
```

```
def go_ignored[X](f:X): Unit
```

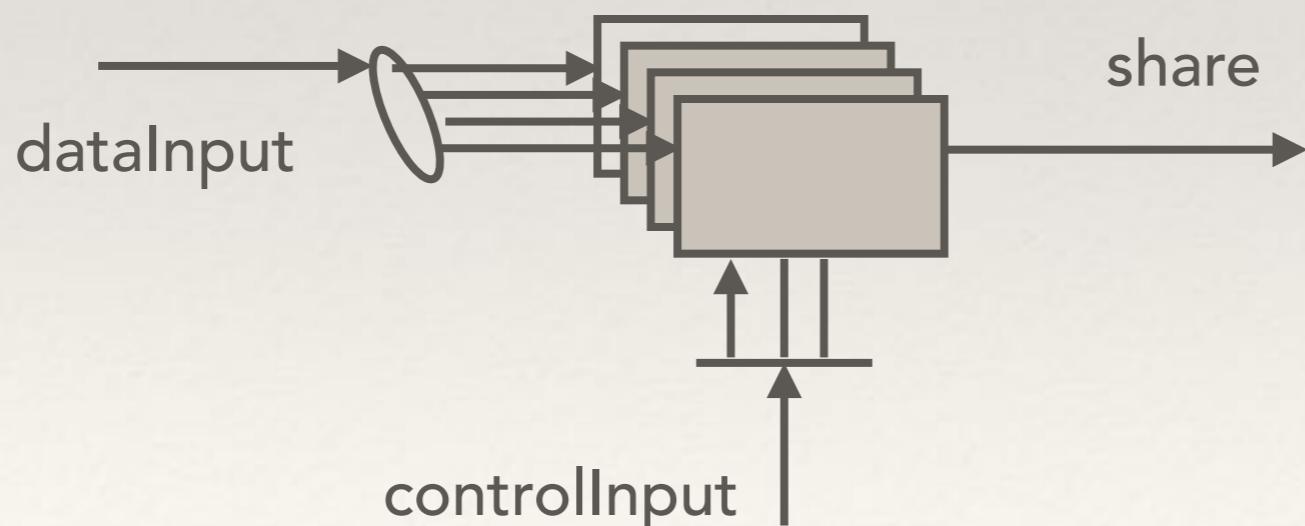
```
def go(f:X): Ignored[Future[X]] =
```

Transputer: select



Transputer: replicate

```
val r = gopherApi.replicate[SMTTransputer](10)
( r.dataInput.distribute( _.hashCode % 10 ) ).
  .controlInput.duplicate().
  out.share()
)
```



Programming techniques

- ❖ Dynamic recursive dataflow schemas
 - ❖ configuration in state
- ❖ Channel-based two-wave generic API
 - ❖ expect channels for reply

Dynamic recursive dataflow

```
select.fold(output){ (out, s) => s match {
  case x:input.read => select.once {
    case x:out.write =>
    case select.timeout => control.distributeBandwidth match {
      case Some(newOut) => newOut.write(x)
        out | newOut
      case None => control.report("Can't increase bandwidth")
        out
    }
  }
  case select.timeout => out match {
    case OrOutput(frs,snd) => snd.close
      frs
    case _ => out
  }
}
```

dynamically increase and decrease bandwidth in dependency from load

Dynamic recursive dataflow

```
select.fold(output){ (out, s) => s match {  
    case x:input.read => select.once {  
        case x:out.write =>  
            case select.timeout =>  
                control.distributeBandwidth match {  
                    case Some(newOut) => newOut.write(x)  
                    out | newOut  
                    case None =>  
                        control.report("Can't increase bandwidth")  
                        out  
                }  
            case _ => out  
    }  
}
```

dynamically increase and decrease bandwidth in dependency from load

Dynamic recursive dataflow

```
select.fold(output){ (out, s) => s match {
  case x:input.read => select.once {
    case x:out.write =>
    case select.timeout => control.distributeBandwidth match {
      case Some(newOut) => newOut.write(x)
        out | newOut
      case None => control.report("Can't increase bandwidth")
        out
    }
  }
}

case select.timeout => out match {
  case OrOutput(frs,snd) => snd.close
    frs
  case _ => out
}
```

dynamically increase and decrease bandwidth in dependency from load

Channel-based generic API

- ❖ Endpoint instead function call
 - ❖ $f: A \Rightarrow B$
 - ❖ endpoint: Channel[A, Channel[B]]
- ❖ Recursive
 - ❖ case class M(A, Channel[M])
 - ❖ $f: (A, M) \Rightarrow M$ (dataflow configured by input)

Channel-based generic API

```
trait Broadcast[T]
{
    val listeners: Output[Channel[T]]
    val messages: Output[T]

    def send(v:T):Unit = { messages.write(v) }

    ....
}
```

- message will received by all listeners

Channel-based generic API

```
class BroadcastImpl[T]
{
    val listeners: Channel[Channel[T]] = makeChannel[Channel[T]]
    val messages: Channel[T] = makeChannel[Channel[T]]

    def send(v:T):Unit = { messages.write(v) }

    select.afold(makeChannel[Message]) { (bus, s) =>
        ....
        s match {
            case v: messages.read => val newBus = makeChannel[Message]
                current.write(Message(newBus,v))
                newBus
            case ch: listeners.read => select.afold(bus) { (current,s) =>
                s match {
                    case msg:current.read => ch.awrite(msg.value)
                        current.write(msg)
                        msg.next
                }
                current
            }
        }
    }
}

// private part
case class Message(next:Channel[Message],value:T)
```

Scala-Gopher: Future directions

- ❖ More experience reports (try to use)
- ❖ Extended set of notifications
 - ❖ channel.close, overflow
- ❖ Distributed case
 - ❖ new channel types with explicit distributed semantics

Scala-Gopher: Conclusion

- ❖ Native integration of CSP into Scala is possible
 - ❖ have a place in a Scala concurrency model zoo
- ❖ Bring well-established techniques to Scala world
 - ❖ (recursive dataflow schemas; channel API)
- ❖ Translation of invocation-only high-order functions into async form can be generally recommended.
 - ❖ (with TASTY transformation inside libraries can be done automatically)

Thanks for attention

- ❖ Questions ?
- ❖ <https://github.com/rssh/scala-gopher>
- ❖ ruslan shevchenko: ruslan@shevchenko.kiev.ua